## **Capacity Analysis of Collaborative Wireless Transmission**

## **Olav Tirkkonen**

Department of Communications and Networking Helsinki University of Technology (TKK) and

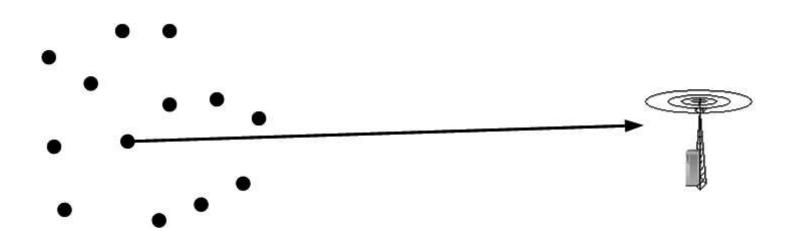
Nokia Research Center, Helsinki

Physics of Distributed Information Systems Workshop Nordita 15.5. 2008



- system model
- random matrices
- results



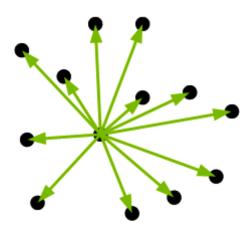


- The source (user) transmits, the destination (BS) receives
- The BS has  $N_{\rm r}$  receive antennas, performs optimum antenna processing
- user at cell edge not too happy



- [Laneman & al 2000, Sendonaris & al 2003, etc, etc]
- clusters of users collaborate to transmit messages
- formation of a virtual antenna array giving rise to matrix (MIMO) channels
- branch of widely investigated relaying field

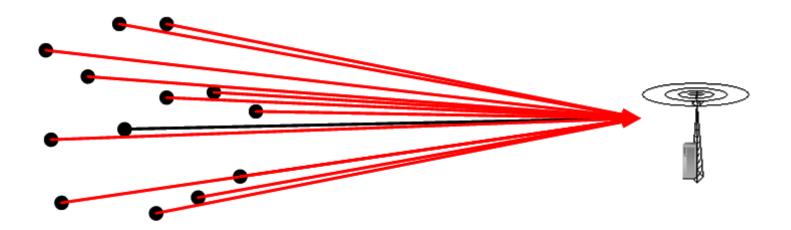






- There are  $N_{\rm t}$  users, one is the source
- The  $N_{\rm t}-1$  other users act as relay nodes
- The source *broadcasts* content to relays
- Assumptions:
  - source knows channel gains  $ilde{\gamma}_j$  between source and relays
  - distance to destination >> distance to relays, direct link in 1:st hop discarded





- two-hop decode and forward relaying
- Relays (and source) transmit content to destination
- Matrix channel, collaborative array transmission from  $N_{\rm t}$  nodes:

$$\mathbf{y}_{\mathbf{v}_{\mathbf{r}} imes 1} \;=\; \mathbf{H}_{N_{\mathbf{r}} imes N_{\mathbf{t}}} \;\; \mathbf{w}_{\mathbf{t} imes 1} \;\; x \;\; + \; \mathbf{n}_{N_{\mathbf{r}} imes 1}$$



• matrix channel SVD:

$$egin{array}{ccc} \mathbf{H} & = & \mathbf{U} & \sum \ _{N_{\mathrm{r}} imes N_{\mathrm{r}}} & \sum \ _{N_{\mathrm{r}} imes N_{\mathrm{t}}} & \mathbf{V}^{\mathrm{H}} \end{array}$$

- collective knowledge of optimum beamforming vector: w maximum eigenvalue vector of  $\mathbf{H}^{\mathrm{H}}\mathbf{H}$ 
  - relays phase and weight transmissions—signals combine coherently at destination, power optimally distributed over channel
  - each relay node knows  $w_j$

• effective channel becomes 
$$\mathbf{H}\mathbf{w} = \mathbf{U}\begin{bmatrix} \sigma_{\max} \\ 0 \\ \vdots \\ 0 \end{bmatrix} = \mathbf{u}_1 \sigma_{\max}$$

- $\sigma_{\max}$  largest singular value of channel
- receiver sees channel gain  $\lambda_{\max} = \sigma_{\max}^2$



- Capacity of hop j is C<sub>j</sub> [nats/s]
- Distribution of  $C_j$  is  $p_j(C_j)$
- Source decides division of resources between 1:st and 2:nd hop:

$$T_1C_1 = T_2C_2$$

Capacity of relay link is

$$C = \frac{T_2 C_2}{T_1 + T_2} = \frac{C_1 C_2}{C_1 + C_2}$$

• distribution of C is

$$p(C) = \int dC_2 \left(\frac{C_2}{C_2 - C}\right)^2 p_1 \left(\frac{CC_2}{C_2 - C}\right) p_2(C_2)$$



• broadcast channel capacity: capacity of worst link

$$C_2 = \min_j \log\left(1 + \tilde{\gamma}_j\right)$$

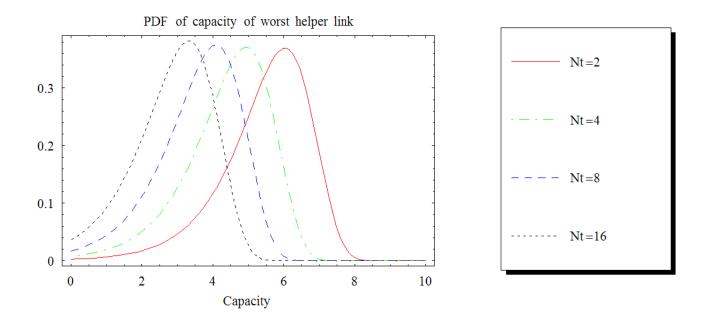
• order statistics, worst of  $N_{\rm t}-1$ 

$$f(\gamma_{\min}) = (N_{t} - 1)F(\gamma)^{(N_{t}-2)}f(\gamma)$$

- channels distributed as  $\sqrt{g_1} \times$  i.i.d complex Gaussian, variance 1:  $\tilde{\gamma}_j = g_1 \gamma_j$  where  $p(\gamma_j) = e^{-\gamma_j}$
- $\Rightarrow$  capacity distribution in closed form

9





 The more friends you want to have, the more likely a bad friend is, the more you have to pay



• H distributed as  $\sqrt{g_2}$  times correlated Gaussian

$$p(\mathbf{H}) = \frac{\mathrm{e}^{-\mathrm{Tr}\left[\mathbf{R}_{r}^{-1}\mathbf{H}\mathbf{H}^{\mathrm{H}}\right]}}{\pi^{N_{\mathrm{t}}N_{\mathrm{r}}} \det^{N_{\mathrm{t}}}\mathbf{R}_{r}}$$
(1)

- correlation only at destination end (realistic)
- assume  $N_{
  m r} \geq N_{
  m t}$ :  ${f H}{f H}^{
  m H}$  is Wishart
- eigendecompose

$$p(\Lambda, U) \sim \frac{\mathrm{e}^{-\mathrm{Tr}\left[\mathbf{R}_{r}^{-1}\mathbf{U}^{\mathrm{H}}\Lambda\mathbf{U}\right]} \mathrm{det}^{N_{\mathrm{r}}-N_{\mathrm{t}}}(\Lambda)\Delta^{2}(\Lambda)}{\mathrm{det}^{N_{\mathrm{t}}}\mathbf{R}_{r}}$$

- now integrate out U
  - an *exactly solvable* partition function on compact phase space
  - Integrability theory: Duistermaat-Heckman can be applied [Morozov 1995]



- for correlation in smaller dimension [Harish-Chandra 1957, Itzykson-Zuber 1992]
- for correlation in larger dimension [Gao-Smith 2000]

• here HCIZ:  

$$p(\Lambda) \sim \frac{\det \mathbf{A} \det^{N_{r}-N_{t}}(\Lambda)\Delta(\Lambda)}{\Delta(\mathbf{R}_{r}^{-1}) \det^{N_{t}} \mathbf{R}_{r}}$$
A is a matrix of exps.  $a = e^{-\lambda_{i}/\rho_{j}}$ 

A is a matrix of exps,  $a_{ij} = e^{-\lambda_i/
ho_j}$ 

•  $\det \mathbf{A}$  carries a contribution from each of the extremal points of the integrand on compact phase space



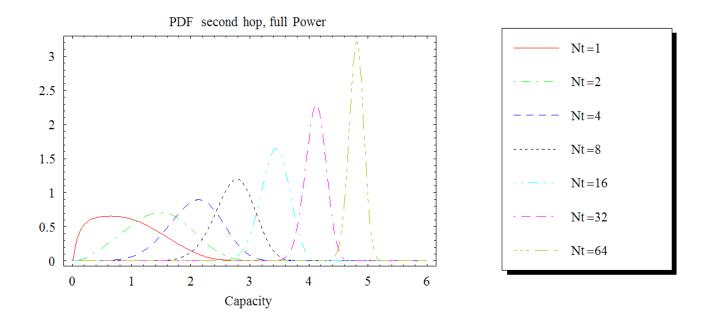
• CDF of alrgest eigenvalue is

$$F(u) = \prod_{j=1}^{N_{\rm t}} \left( \int_0^u \mathrm{d}\lambda_j \right) p(\Lambda)$$

- [Dighe & al 2003] integrated non-correlated case, use same trick
- $\det A$  is a sum (over permutations of correlation eigenvalues) of integration measures. Treat one-by-one.
- expression to integrate is  $\det \mathbf{E} \equiv \det \left( \Lambda^{N_{r}-N_{t}} \mathbf{D}(\Lambda) \right)$ 
  - $\mathbf{D}(\Lambda)$  is the VanDermonde matrix,  $d_{mn} = \lambda_m^{n-1}$ .
  - elements of  $\mathbf{E}$  are thus  $e_{mn} = \lambda_m^{N_{\mathrm{r}} N_{\mathrm{t}} + n 1}$ .
  - each row of  ${f E}$  depends only on one  $\lambda_j$
  - integration decouples to product of trivial 1D integrals

 $\Rightarrow \lambda_{\max}$  distribution in closed form





- $C_2 = \log\left(1 + g_2 \lambda_{\max}\right)$
- closed form distriburtion  $p_2(C_2)$
- The more friends you have, the more they may help you
- with numerous friends, mean-field bahavior

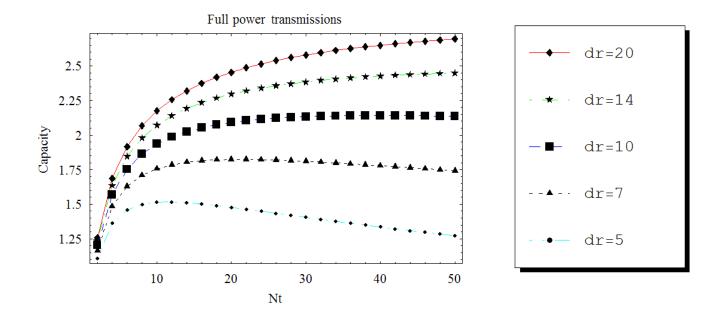


- $g_2 = 1$  (Signal-to-Noise Ratio on relay link)
  - Relay links are weak, cluster far from destination
- same power used by the source in the first hop, and on the average by the source and relays in the second hop
  - the more relay nodes, the more power radiated into air

• 
$$g_2 = d_r^{\alpha}/G_r$$

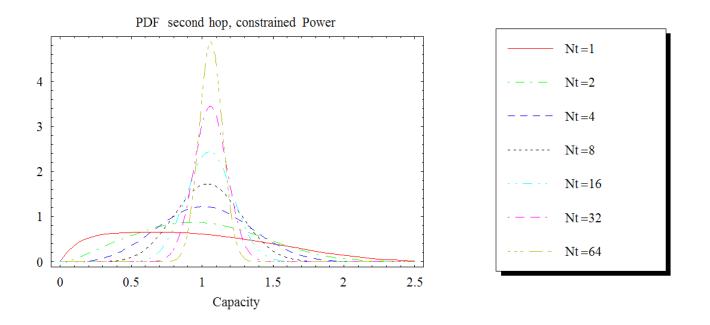
- Path-loss exponent  $\alpha = 3.75$
- destination antenna gain  $G_r = 10$
- $d_r$  ratio between radius of cluster and distance b/w cluster and destination





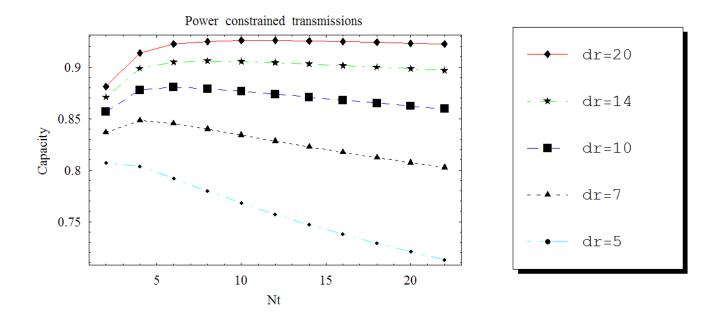
- without relaying, the capacity is 0.94
- cost vs benefit tradeoff
- the furhter away you search for friends, the more you pay





- Strict control on interference: constrained total power usage in cell
- some array gain visible (optimum usage of differnet channels)





- without relaying, the capacity is 0.94
- The cost outweights the gains
- outage analysis (reliability) would look different



- Analyzed cost-benefit tradeoff on collaborative relaying for cell-edge users
- No gain if not allowed to radiate more power into cell
- this would improve with instantaneous selection of relay cluster
- gains come at a cost of the batteries of friends
- allowing increased power usage, significant iprovements to cell-edge service