

# Uniform sampling of local minima in Ising lattices

Pekka Orponen  
joint work with Petri Savola

Department of Information and Computer Science  
Helsinki University of Technology TKK

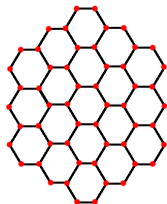
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## Motivation, background

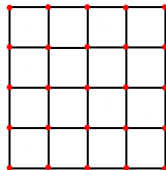
- ▶ For exploring the energy landscape of spin glass models, it is useful to have an efficient method for sampling local energy minima (metastable states).
- ▶ MCMC methods can be adapted to this purpose.
- ▶ We present a direct combinatorial method based on a classical algorithm by Knuth (1975) for estimating the size of large search trees.
- ▶ Applications: estimating the number of minima as a function of system size; exploring the energy distribution of minima; exploring the attraction basins around minima at various energy levels; ...

## Models

We consider the following two types of 2D Ising spin glass lattices:



(a) Hex lattice

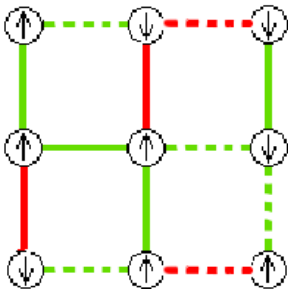


(b) Square lattice

- ▶ Up-down spins:  $\sigma_i \in \{1, -1\}$
- ▶ Positive-negative interactions:  $J_{ij} \in \{1, -1, 0\}$
- ▶ Hamiltonian:  $H(\sigma) = -\frac{1}{2} \sum_{ij}^N J_{ij} \sigma_i \sigma_j$

## Graphical view of local minima (1/2)

- ▶ A spin configuration  $\sigma$  induces a colouring of the edges in the respective lattice into “green” (satisfied) and “red” (unsatisfied)

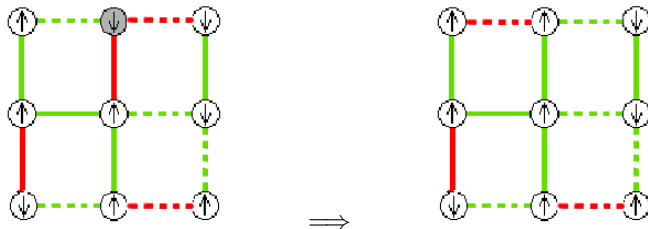


Here solid edge indicates a positive interaction, a dotted edge a negative one.

- ▶ *Note:*  $H = -(\# \text{ of green edges}) + (\# \text{ of red edges})$

## Graphical view of local minima (2/2)

- ▶ An edge colouring of the lattice corresponds to a (proper) local energy minimum, if and only if each vertex is incident to (properly) fewer red than green edges.
- ▶ E.g. the previous example colouring (spin configuration) can still be improved by flipping the top spin:



- ▶ Technically, there is also another validity condition on a colouring: the induced spin orientations around any cycle in the lattice must be consistent. (E.g. in a 4-cycle of positive interactions, there must be an even number of red edges.)

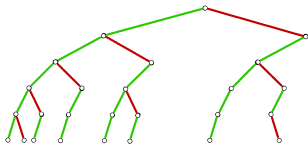
## Enumerating local minima (1/2)

- ▶ Thus, a local minimum corresponds to an edge colouring of the lattice (interaction graph)  $G$  satisfying a local consistency condition. (Because of the up-down symmetry of the spins, there are actually two minima per each valid colouring.)
- ▶ Consider any spanning tree  $T$  of  $G$ .
- ▶ Any edge colouring of  $G$  obviously yields a unique colouring of  $T$ .
- ▶ More interestingly, any colouring of  $T$  has a unique completion to all of  $G$ . (Propagate the spin orientations from the root of  $T$  according to the colour of the tree edges. Colour the non-tree edges of  $G$  accordingly.)
- ▶ Starting from the root vertex, the local validity conditions are easy to satisfy on  $T$ . (But the completion to  $G$  does not necessarily satisfy them.)

## Enumerating local minima (2/2)

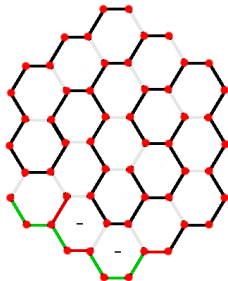
An algorithm to enumerate local minima on an interaction graph  $G$ :

- ▶ Choose spanning tree  $T$  of  $G$ . List edges of  $T$  in some order.
- ▶ Enumerate systematically all colourings of this edge list ( $\implies$  binary search tree).
- ▶ Only consider branches of the search tree where the current partial colouring can be completed to a valid colouring of all of  $G$ . (This may be difficult to test!)
- ▶ Now leaves of the search tree correspond one-to-one to consistent colourings ( $\sim$  local minima) of  $G$ .



## Complexity of completability

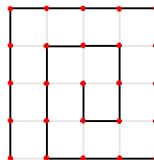
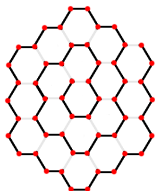
- ▶ In a general spanning tree, determining the completability of a given partial colouring may require arbitrary long look-ahead:





## Spanning paths

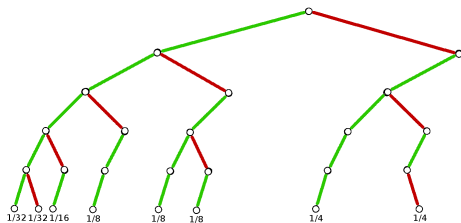
- ▶ In lattices, spiral-like paths provide particularly simple spanning trees:



- ▶ Moreover, if the edges of the spanning spiral are coloured inside-out, the completability can be tested efficiently. (Basically, a partial colouring can always be completed, unless the validity conditions are violated by already chosen or “locally forced” edge colours.)

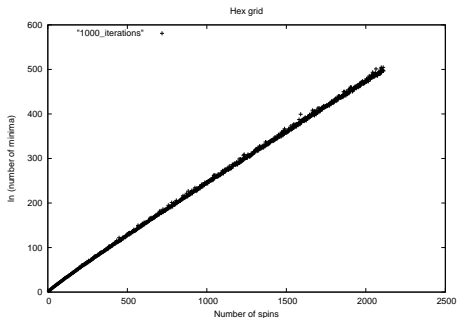
# Estimating the size of a search tree

- ▶ D. E. Knuth, “Estimating the efficiency of backtrack programs”. Math. Computation 29 (1975), 121–136.
- ▶ Method to estimate size = number of leaves  $S$  in a large search tree  $T$ :
  - ▶ Make a random descent into  $T$ , starting from the root.
  - ▶ Record degrees (branching factors)  $d_1, d_2, \dots, d_n$  of the vertices encountered along the descent. (For a binary tree,  $d_i \in \{1, 2\}$  for each  $i$ .)
  - ▶ Compute estimate  $\hat{S} = d_1 d_2 \cdots d_n$ .
- ▶ *Theorem.*  $\hat{S}$  is an unbiased estimate of the true tree size  $S$ .
- ▶ To decrease variance of the estimate, perform several descents.



# Application 1: Estimating the number of minima (1/2)

- ▶ Ansatz:  $S \sim e^{\alpha N}$ , where  $S$  is the number of local minima,  $N$  is the number of spins, and  $\alpha$  is a coefficient which depends on the system characteristics.
- ▶ E.g. for hexagonal lattices of up to 2000 spins, with 50% fraction of negative interactions, we obtain the following diagram. (Each estimate of  $L$  is based on 1000 descents.)



## Application 1: Estimating the number of minima (2/2)

- ▶ Numerical estimates of the coefficient  $\alpha$ , from systems of 1 ... 2000 spins:
  - ▶ Ferromagnetic hex lattice:  $\alpha \approx 0.226$
  - ▶ Hex lattice with 50% negative interactions:  $\alpha \approx 0.231$
- ▶ The results match quite well the analytic prediction of  $\alpha \approx 0.225$  from a recent paper by Waclaw and Burda (arXiv Jan 2008). The paper also reports experimental data on systems of up to 24 spins, indicating  $\alpha \approx 0.226$ .

## Uniform sampling of leaves in a search tree

- ▶ Apply Knuth's method recursively.
- ▶ When at a non-leaf vertex of a binary search tree  $T$ :
  - ▶ Estimate size of left subtree:  $\hat{S}_L$
  - ▶ Estimate size of right subtree:  $\hat{S}_R$
  - ▶ Descend to left subtree with probability

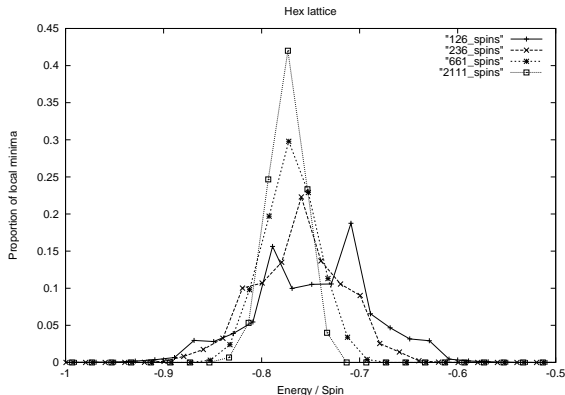
$$p = \frac{\hat{S}_L}{\hat{S}_L + \hat{S}_R},$$

and to right subtree with probability  $1 - p$ .

⇒ Descent reaches each leaf with equal probability.

## Application 2: Energy distribution of local minima

- Based on a uniform sampling of 1000 minima in ferromagnetic hex lattices with  $N = 126, \dots, 2111$  spins. (Equidistant binning of energy levels. Only one descent per subtree size estimate.)



## Summary and future work

- ▶ New(?) combinatorial method for uniform sampling of local minima in spin glass lattices.
- ▶ Method has potentially many applications to the exploration of spin glass energy landscapes.
- ▶ Basic method scales well to large system sizes, but effect of system size and search tree shape on the variance of the method are not well understood.
- ▶ Future work 1: Analyse variance of the method.
- ▶ Future work 2: Extend the method from 2D lattices to other graph structures.
- ▶ Future work 3: Explore other applications in landscape analysis.