

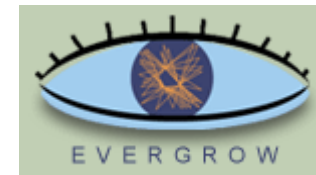
Statistical Mechanics of LDPC Codes in MIMO Channels

(<http://arxiv.org/abs/0705.1644v1>)

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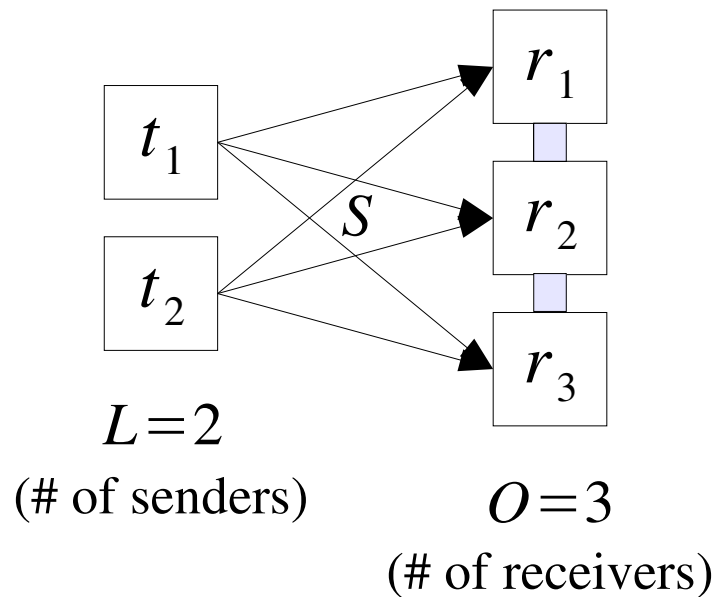
Statistical Mechanics of Distributed Information Systems
15-18 July 2007, Mariehamn, Finland

Summary

- Multiple-Input Multiple-Output Channels
- Low-Density Parity-Check Codes
- Decoding: MPM estimator
- Statistical Physics Formulation: Replica Symmetric Ansatz
- Channels:
 - Single Transmitter
 - Multiple Access
 - Interference Channel (Symmetric, Asymmetric)
- Conclusions

Multiple-Input Multiple-Output (MIMO) Channels

Joint Decoding



Transmitted codewords

Additive White
Gaussian Noise

$$\vec{r}^\mu = S \vec{t}^\mu + \vec{v}^\mu$$

Received vectors

Interference Matrix

Low-Density Parity-Check (LDPC) Codes

Generator Matrix (G): Original Message (w) --> Codeword (t)

$$G = \begin{bmatrix} I \\ B \end{bmatrix} \Rightarrow t = \begin{bmatrix} w \\ Bw \end{bmatrix} \quad \longrightarrow \quad w \rightarrow t = Gw \pmod{2}$$

Decoding

Parity Check Matrix

$$AG = 0 \Rightarrow At = AGw = 0$$

Syndrome

$$r = t + v \Rightarrow Ar = Av \equiv z$$

Gallager Codes

$$A = [C_1 | C_2]$$

$$G^T = [I | C_2^{-1} C_1]$$

Regular Codes: $R = 1 - \frac{C}{K}$

*nonzero entries
per column*

*nonzero entries
per row*

Decoding: **M**arginal **P**osterior **M**aximiser (**MPM**)

Encoding: i.i.d. Parity-Check Matrices for each sender.

MPM Decoding

$$\hat{t}_i^\mu = \text{sgn} \langle \tau_i^\mu \rangle_{P(\tau|r)}$$

minimizes bit error

Overlap

$$d_i = \frac{1}{M} \sum_{\mu=1}^M \langle t^\mu \hat{t}_i^\mu \rangle_{A_1, \dots, A_L, r, t}$$

Fast Decoding by
Probability (Belief) Propagation

Typical behaviour averaged over
all codes, codewords and noise.

Statistical Physics Formulation: Replica Symmetric Ansatz

$$P(\tau|r) = \frac{P(r|\tau)P(\tau)}{P(r)} \rightarrow P(\tau|r) = \frac{e^{-\beta H(\tau|r)}}{Z}$$

inference problem

SM problem

$$H(\tau|r) = -\ln [P(r|\tau)P(\tau)]$$

Nishimori's Temperature

$$\beta = 1$$

Example: Parity Check Code + Gaussian Channel

$$r = \tau + \nu$$

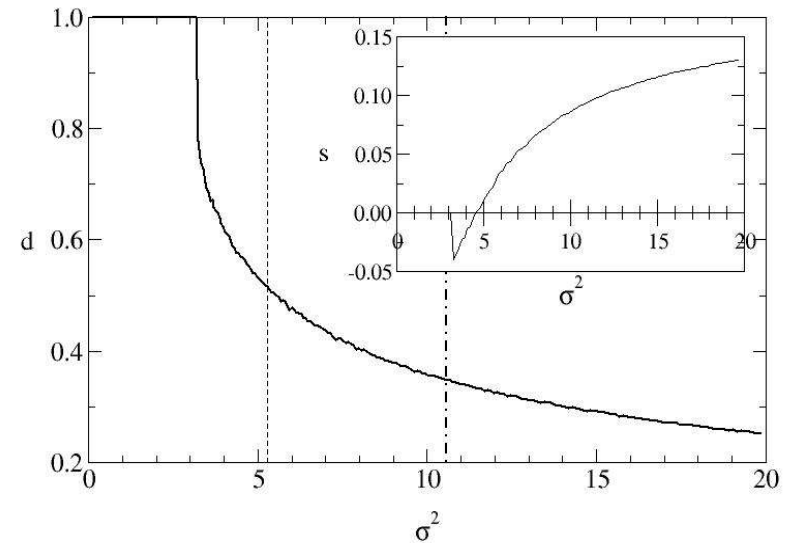
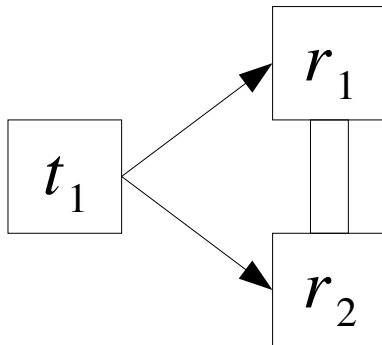
$$P(r|\tau) \propto e^{-\frac{1}{2\sigma^2}(r-\tau)^2}$$

local field

$$P(\tau) \propto \chi(A\tau = 0)$$

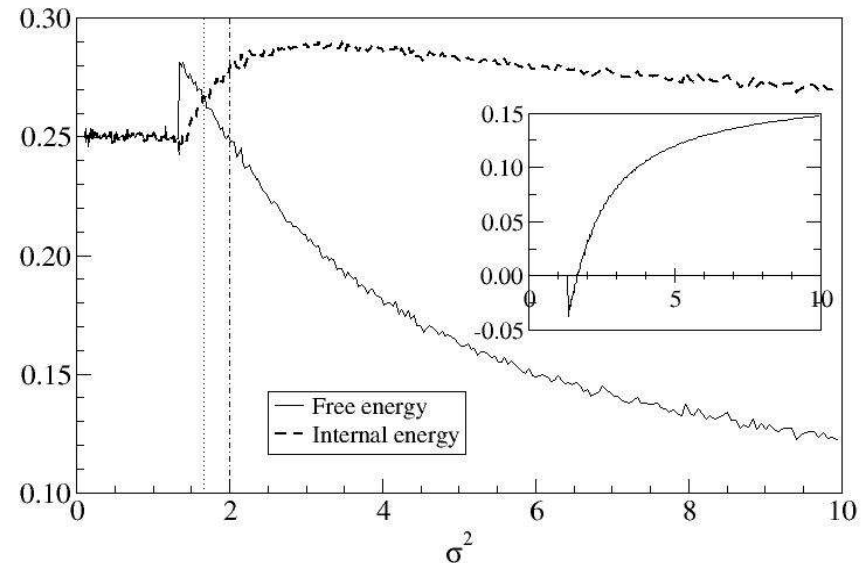
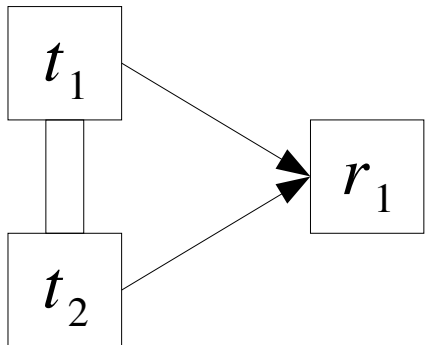
K-spin
interaction

Single Transmitter



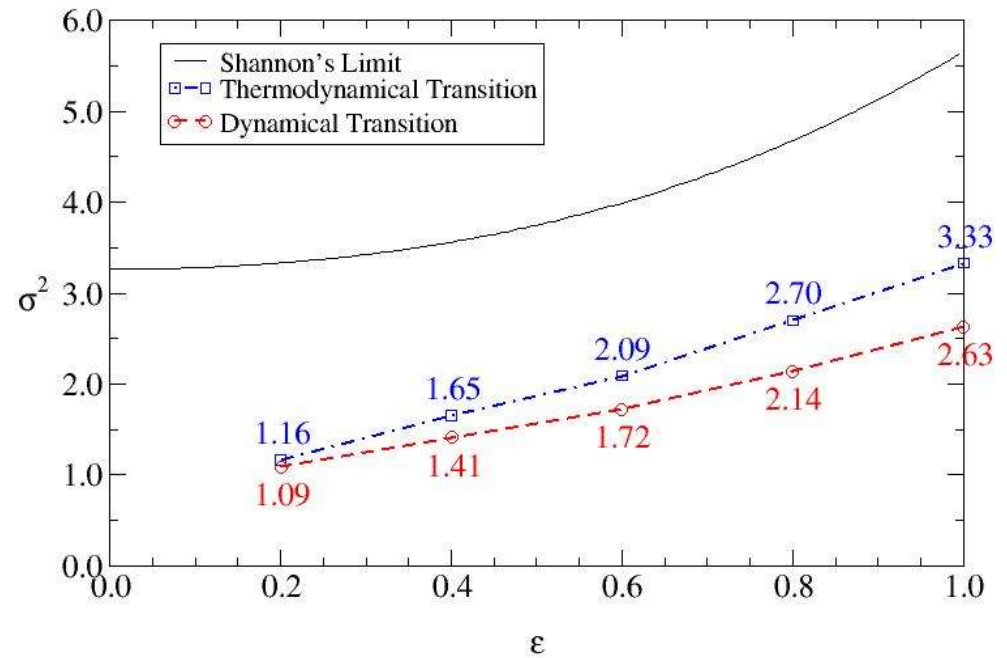
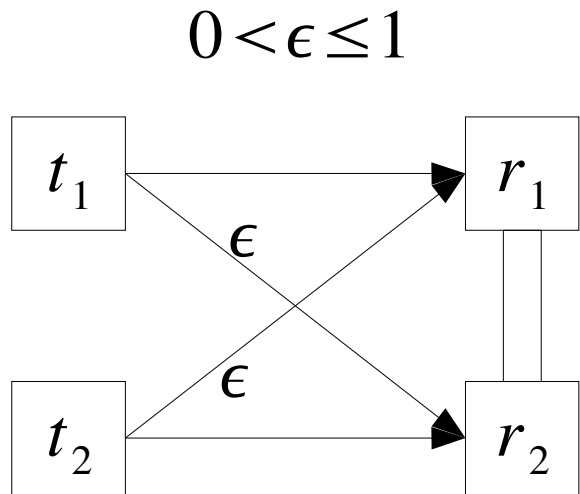
O	Shannon's Limit	Dynamical Transition	Thermodynamical Transition	T.T.-D.T.	S.L. Gaussian Channel	O messages
1	2.41	1.59	2.24	0.65	2.41	
2	10.57	3.28	4.59	1.31	5.28	
3	24.50	4.90	6.68	1.78	8.17	

Multiple Access Channel (MAC)

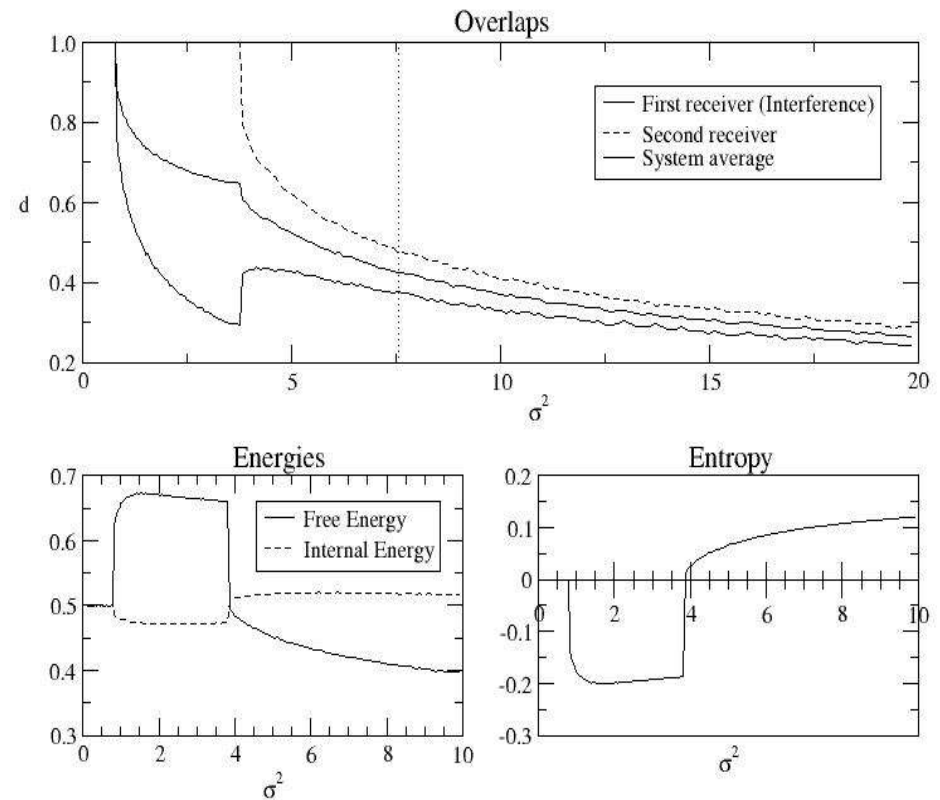
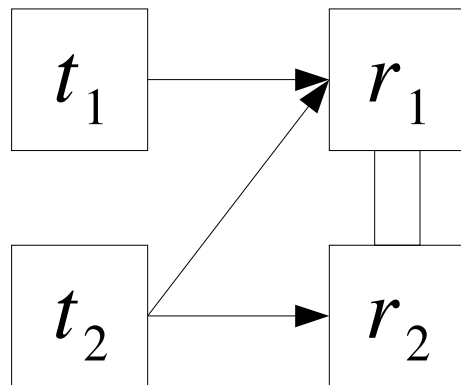


\mathcal{O}	Shannon's Limit	Dynamical Transition	Thermodynamical Transition	T.T.-D.T.
1	2.41	1.59	2.24	0.65
2	2.00	1.32	1.66	0.34
3	1.64	1.24	1.45	0.21

Symmetric Interference Channel



Asymmetric Interference Channel



Conclusions

- Improvement/deterioration in MIMO channels
- Role of metastable states generated by L and O
- Importance of Network Coding
- Increasing threshold with increasing interference in the symmetric interference channel
- Non-trivial information sharing in asymmetric channels