

# Quantum annealing and glass problems

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**Quantum computers are: a) good for Fourier Transform-based algorithms** coherence

b) (obviously) good for avoiding the sign problem

no-coherence ?

c) most probably *bad* for "Golf Course"  
(=Random Energy=Glassy) systems

# Sign Problem

**We are asked to compute averages  
over configurations  $s$**

$$\langle O \rangle = \frac{\sum_s e^{-W(s)} O(s)}{\sum_s e^{-W(s)}}$$

**and  $W = W_R + iW_I$  is not real.**

- **Quantum mechanics / Field Theory with real time or  $\theta$ -terms**
- **Hubbard-Stratonovich decoupling**  $W = W_o - b \sum C_\alpha^2$

$$Z = \sum_s e^{-W_o(s) - b \sum C_\alpha^2} = \sum_s \int d\lambda e^{-W_o(s) + \sqrt{b} \lambda_\alpha C_\alpha - \frac{\lambda_\alpha^2}{2}}$$

The Hubbard-Stratonovich transformation above introduces an imaginary term  $W_I = \sqrt{|b|} \sum \lambda_\alpha C_\alpha$  in the repulsive case  $b < 0$ .

**In solid state theory the sign problem is the main obstacle for giving a numerical answer to very urgent questions.**

e.g. Hubbard model  $H = -t \sum_{\langle ij \rangle, \sigma} (c_{i\sigma}^\dagger c_{j\sigma} + h.c.) + U \sum_i n_{i\uparrow} n_{i\downarrow}$

**For definiteness:**  $W_I = ih_I M(s)$

where  $M(s)$  is an integer-valued:

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$$Z = \sum_M Z_M e^{-ih_I M}$$

$$Z_M = e^{-\beta F(M)} = \sum_s \delta(M(s) - M) e^{-W_R}$$

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- **Field theory with  $\theta$  terms.**  $M$  = a topological number 't Hooft, Haldane,...
- **Fermion systems:**  $M$  = Fermion world-line crossings Muramatsu et al
- **Hubbard model:**  $M$  = the number of up Hirsch spins

[with a variant of Hubbard-Stratonovich transformation, see: DeForcrand, Batrouni]



**Also: If we knew how to deal with the sign problem, we would know how to compute numerically averaged disordered models**

**using identities of the form**

$$\frac{1}{x} = -\lim_{\lambda \rightarrow \infty} \frac{1 - e^{-\lambda x}}{x} = -\lim_{\lambda \rightarrow \infty} \sum_1^{\infty} (-1)^n \frac{\lambda^n}{n!} x^{n-1}$$

—————

$$\overline{\langle E \rangle} = -\overline{Z^{-1} \frac{\partial Z}{\partial \beta}} = -\lim_{\lambda \rightarrow \infty} \frac{\partial}{\partial \beta} \sum_1^{\infty} \frac{(-1)^n}{n} \frac{\lambda^n}{n!} \overline{Z^n}$$

**In practice, what we can do is to compute**

$$\langle \mathbf{O} \rangle = \frac{\langle e^{-i\mathbf{W}_I} \mathbf{O} \rangle_R}{\langle e^{-i\mathbf{W}_I} \rangle_R} \quad ; \quad \text{where} \quad \langle \bullet \rangle_R = \frac{\sum \bullet e^{-W_R(s)}}{\sum e^{-W_R(s)}}$$

Note that Monte Carlo is only really good to calculate  $\langle O \rangle$  for  $O$  non-exponential

**and not for**

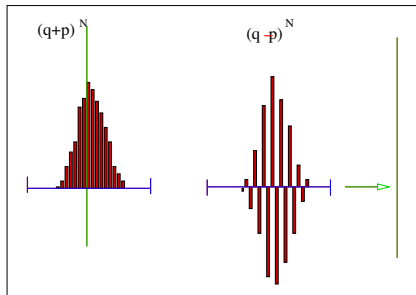
$$\langle O \rangle = \frac{\langle e^{-B} O \rangle_A}{\langle e^{-B} \rangle_A}$$

**if, e.g.  $B$  is exponential in  $N$**

## An example:

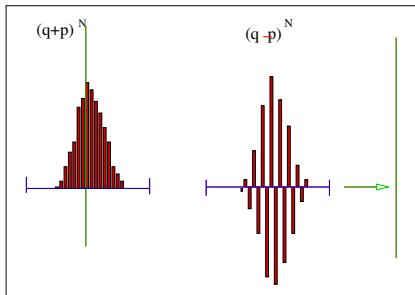
Non-interacting spins in a magnetic field  $h_R + i\frac{\pi}{2}$

$$Z = \left( e^{-h-i\frac{\pi}{2}} + e^{h+i\frac{\pi}{2}} \right)^N = e^{-i\frac{N\pi}{2}} \left( e^{-h} - e^h \right)^N$$



$$(q + p)^N = \sum_r \frac{N!}{r!(N-r)!} q^{N-r} p^r$$

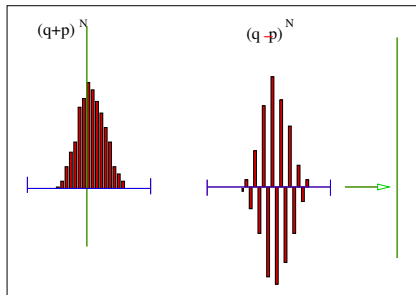
$$(q - p)^N = \sum_r \frac{N!}{r!(N-r)!} (-1)^r q^{N-r} p^r$$



put  $x = \frac{r}{N}$  and use Stirling:

$$(q + p)^N = \int_0^1 dx \, e^{N[-x \ln x - (1-x) \ln(1-x) + (1-x) \ln q + x \ln p]}$$

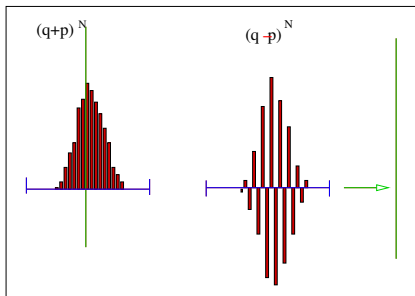
$$\text{saddle: } \left( \frac{x^*}{1-x^*} \right) = \left( \frac{p}{q} \right) \in [0, 1]$$



put  $x = \frac{r}{N}$  and use Stirling:

$$(q - p)^N = \int_0^1 dx \, e^{N[-x \ln x - (1-x) \ln(1-x) + (1-x) \ln q + x \ln p + i\pi x]}$$

saddle:  $\left( \frac{x^*}{1-x^*} \right) = - \left( \frac{p}{q} \right) \quad \text{not} \in [0, 1]$



## Monte Carlo Sampling of the same problem:

$$(q - p)^N \sim \sum_r \left\{ \frac{N!}{r!(N-r)!} + \sqrt{\frac{N!}{r!(N-r)!}} \eta(r) \right\} \quad (-1)^r q^{N-r} p^r =$$

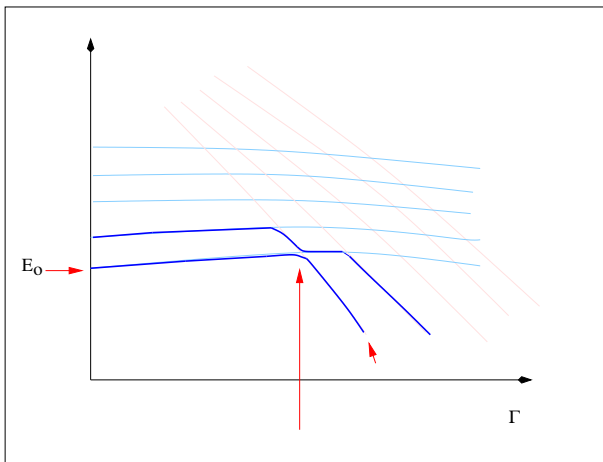
$$(q - p)^N \sim \sum_r \left\{ e^{NS(r)} + e^{N \frac{S(r)}{2}} \eta(r) \right\} \quad (-1)^r q^{N-r} p^r \quad \text{who wins?}$$

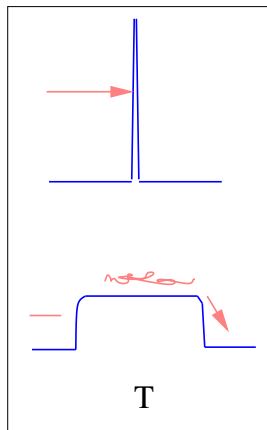
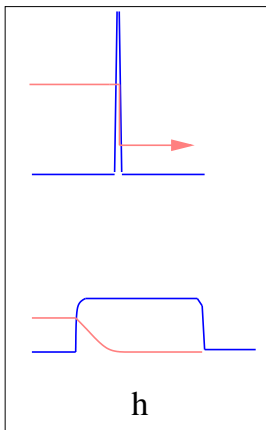
# Quantum Annealing of Hard Problems



$$\mathcal{H}(\{\sigma\}) = E(\{\sigma^z\}) + \Gamma \sum_{i=1}^N \sigma_i^x = \mathcal{H}_0 + \Gamma V \quad (1)$$

Staying in the lowest level without de-railing requires speed  $\Delta^{-2}$



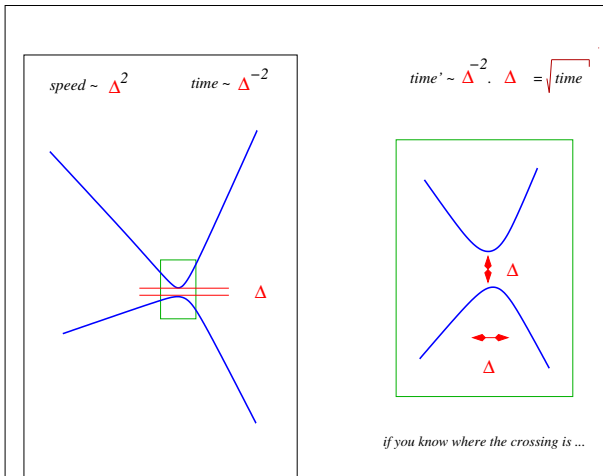


Fat and small versus thin and tall...

The speed is directly given by the minimal gap.

Adiabatic Quantum Computation is Equivalent to Standard Quantum  
Computation (Aharonov et al)

If you know where the crossing takes place, you gain a square root



# Grover

Finding a needle in a haystack, if you know its color...

$$\begin{bmatrix} 1 & & & & \\ & 0 & & & \\ & & 0 & & \\ & & & -1 & \\ & & & & 0 \end{bmatrix} - \Gamma \begin{bmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & \dots & \dots & \dots & 1 \\ & & & & 1 \\ & & & & 1 \\ 1 & 1 & \dots & \dots & 1 \end{bmatrix}$$

WHERE IS THE -1 ?

From  $N$  to  $\sqrt{N}$ , non-trivial yet non miraculous...

# The Random Energy Model.

Energies are independent Gaussian random numbers.

An idealization of the  $p$  ( $> 2$ )-spin model...

$$E(\{\sigma^z\}) = \lim_{p \rightarrow \infty} \sum_{i_1, \dots, i_p} J_{i_1, \dots, i_p} \sigma_{i_1}^z \dots \sigma_{i_p}^z$$

K-SAT...

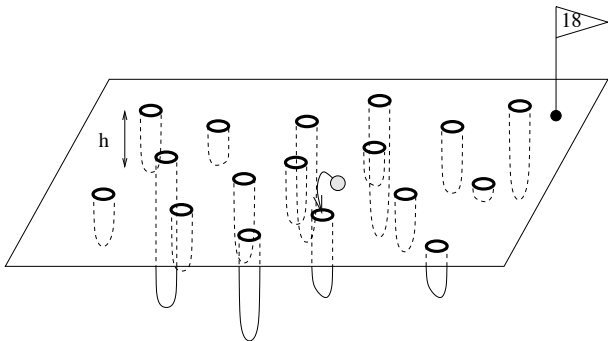
$$E(\{\sigma^z\}) = \sum_{\text{clause } a} C_a \quad C_a = 0 \text{ iff the clause } a \text{ is satisfied.}$$

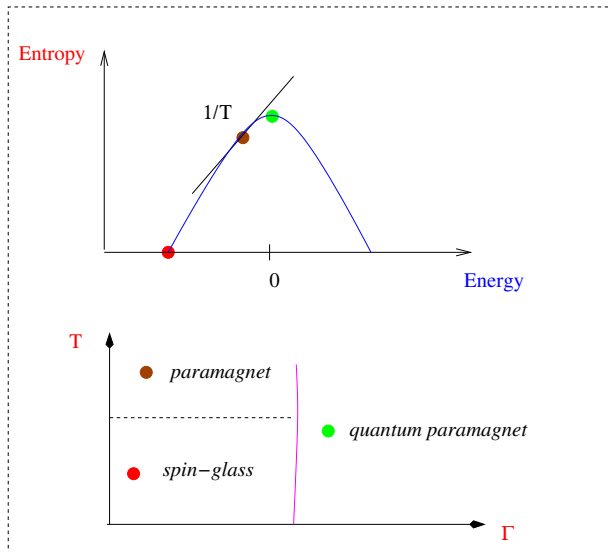
and many other glass models.

# The Random Energy Model.

Energies are independent Gaussian random numbers.

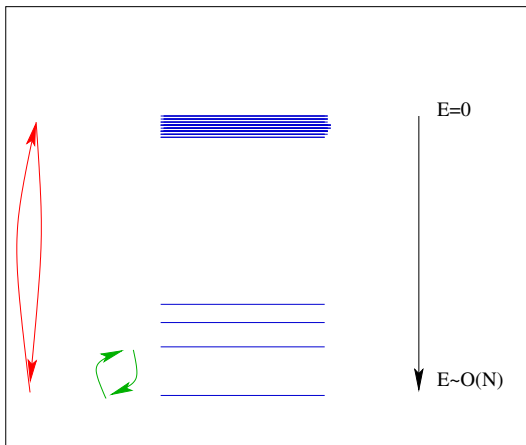
each basin is schematised by a single configuration

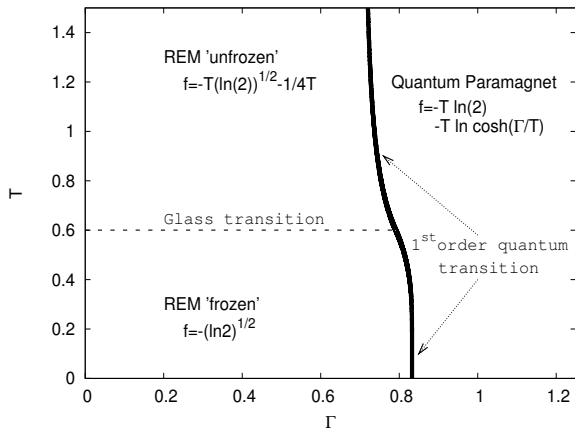


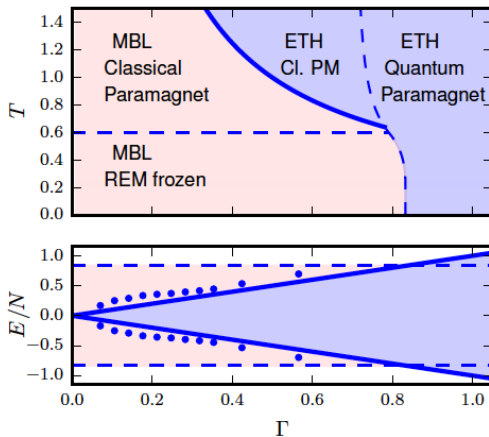




$$E_i(\Gamma) = E_i + \Gamma V_{ii} + \sum_{k \neq i} \frac{\Gamma^2 V_{ik} V_{ki}}{E_i(\Gamma) - E_k} + \dots = E_i + \frac{N\Gamma^2}{E_i} + O\left(\frac{1}{N}\right)$$





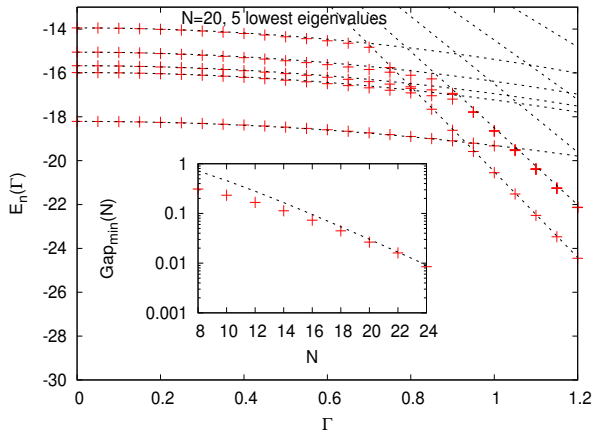


To compute the gap, we just have to diagonalise:

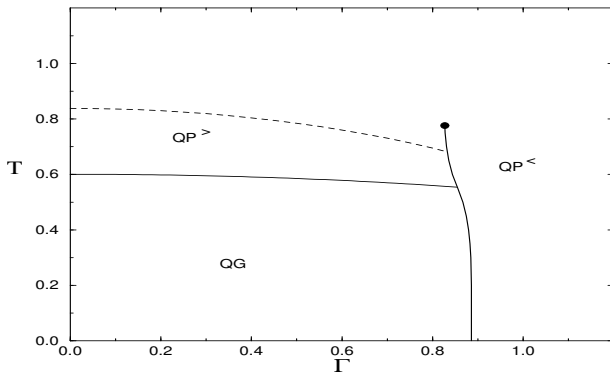
$$\mathcal{H}|\phi\rangle = [E_o|SG\rangle\langle SG| - \Gamma N|QP\rangle\langle QP|]|\phi\rangle = \lambda|\phi\rangle$$

The gap is exponentially small

$$\Delta_{\min}(N) = 2|E_o|2^{-N/2}$$



# Generic Random First Order random $p > 2$ -spin, Potts, etc etc



## Suzuki-Trotter + Replica Approach

Order parameter:  $q^{\mu\nu}(t, t')$

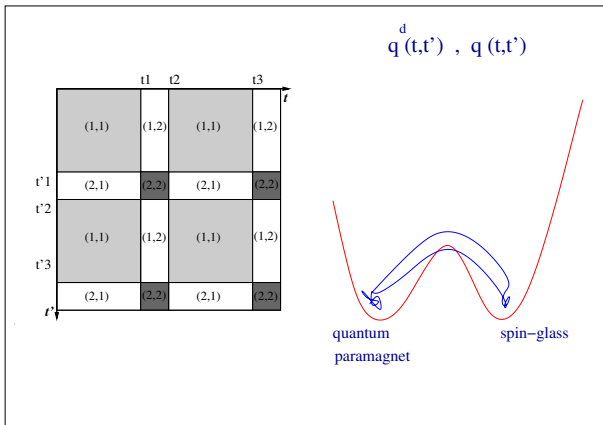
A replicated closed polymer in a random potential!

## One-step RSB ansatz

Diagram illustrating the structure of the matrix  $q_{tt'}^{\mu\mu'}$ . The matrix is shown as a block matrix with square blocks on the diagonal and zero blocks off-diagonal. The top-left block is labeled with  $q_{tt'}^d$  and  $q_{tt'}$ . A double-headed arrow at the bottom indicates the size of the blocks is  $m$ .



# A two-time instanton Approach



In general, the Gap is (minus) the exponential of the free-energy cost of a two-time wall

It is hence the exponential of a negative **extensive** quantity

One can easily recover the result of the REM

**In conclusion, this class of hard problems  
remains exponentially hard in Quantum  
Annealing**

**Not surprising: you do not thermalize a glass in  
real life by making it quantum.**

## Sign problem is more mysterious

**it is understandably not hard for a quantum computer, but does not seem to require coherence (?)**