

Constructing the 6 Loop 4 Point $\mathcal{N} = 4$ sYM Integrand

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with

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Roadmap

- 1 Motivations and Introductions
- 2 Uncrossing (X) and Hiding (H) Identities
- 3 Uncrossing and Hiding in Method of Maximal Cuts (MMC)
- 4 Six Loop Construction

Motivations

Goal of program: UV behavior of 7 loop $\mathcal{N} = 8$ SUGRA

Why?

- SUSY arguments predict $L = 7$ counterterm in $D_C = 4$ (Bossard, Howe, Stelle; Green, Russo, Vanhove; Beisert, Elvang, Freedman, Kiermaier, Morales, Stieberger; many more)
- Similar counterterms proven absent for $\mathcal{N} = 4, 5$ at $L = \mathcal{N} - 1$ (Bossard, Howe, Stelle, Vanhove; Bern, Davies, Dennen; Bern, Davies, Dennen, Huang)
- Improved behavior observed in $D = 4$ kinematics (AE, Hermann, Parra-Martinez, Trnka)

History of direct calculations:

- 1&2 loops '80 - '90s (Green, Schwarz, Brink; Bern, Dixon, Dunbar, Perelstein, Rozowsky)
- 3 loops '07-'10 (Bern, Carrasco, Dixon, Johansson, Kosower, Roiban)
- 4 loops '09-'12 (Bern, Carrasco, Dixon, Johansson, Roiban)
- 5 loops 2018 (Bern, Carrasco, Chen, AE, Johansson, Parra-Martinez, Roiban, Zeng)

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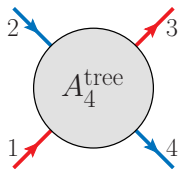
Computational Challenges

Obstacle	Solution
Many Feynman diagrams, cancellations between diagrams	Cuts contain minimal needed data
Cuts from state sums: 256 states per cut propagator	Double copy: $GR = YM^2$
sYM state sums still hard, rapidly exploding γ traces	Color-kinematics + tricks for special cuts
Problems with CK at 5L – still need to solve previous problems	New recursive tools for cuts & integrands

X Identity

Four-point ordered YM tree amplitudes only have s and t channel poles.

What if we try to “sit on the u pole” anyway, via $p_3 \rightarrow p_1, p_4 \rightarrow p_2$?



e.g.:

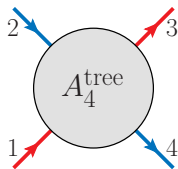
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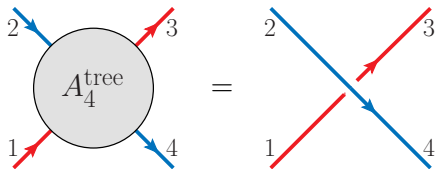
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What if we try to “sit on the u pole” anyway, via $p_3 \rightarrow p_1, p_4 \rightarrow p_2$?

Dimensionless, must respect all symmetries: Can only get identity insertions! The diagram disconnects!



e.g.:

$$A_4^{\text{tree, YM}} = \frac{t_8 F^4}{s t} \xrightarrow{p_3=p_1} (\epsilon_1 \cdot \epsilon_3)(\epsilon_2 \cdot \epsilon_4)$$

Same for all other supersymmetric states

H Identity

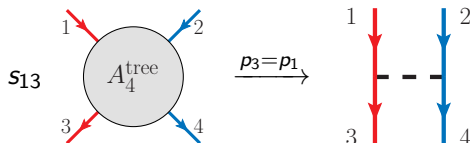
Can we find a similar identity that maintains planarity?

- Consider $s_{13}A(1, 3, 2, 4)|_{p_3=p_1}$: cut with zero momentum exchange
- Apply $s_{13}A(1, 3, 2, 4) = -s_{12}A(1, 2, 3, 4)$
- Use X ID on RHS: $s_{12}A(1, 2, 3, 4) \rightarrow -s_{14}A(1, 3)A(2, 4)$

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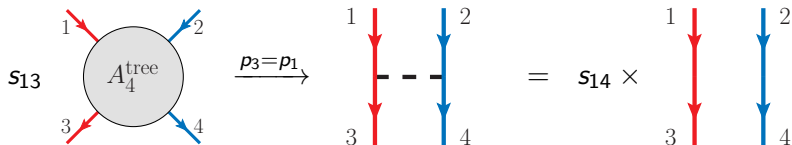
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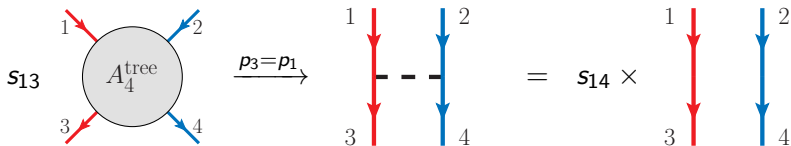
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$$\text{e.g.: } s_{13}A(1, 3, 2, 4) = \frac{t_8 F^4}{s_{23}} \xrightarrow{p_3=p_1} s_{14}(\epsilon_1 \cdot \epsilon_3)(\epsilon_2 \cdot \epsilon_4)$$

Same for all other supersymmetric states. Extends to (super)gravity.
 N.B.: Physical identity (soft factorization), not heuristic rule.

Method of Maximal Cuts (MMC)

(Bern et al)

Systematic integrand construction from cuts

- ① Enumerate diagram basis, striate by *cut depth* (k)
 - ① Build all cubic vacuums
 - ② Attach four external legs
 - ③ Collapse internal legs
 - ④ Cubic = max; one quartic = next-to-max; ...
- ② Each diagram γ corresponds to both a **cut** and a **numerator**
- ③ Proceed by cut level, matching cuts by inheriting poles from lower- k and constructing new numerators

$$\mathcal{P}_{\gamma,\text{ans}}^{(k)} = \mathcal{C}_{\gamma}^{(k)} - \mathcal{R}_{\gamma,\text{MMC}}^{(k)}$$

Local polynomial
Cut \rightarrow truth
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The diagram illustrates the equation $\mathcal{P}_{\gamma,\text{ans}}^{(k)} = \mathcal{C}_{\gamma}^{(k)} - \mathcal{R}_{\gamma,\text{MMC}}^{(k)}$. Arrows point from each term to a label below it: $\mathcal{P}_{\gamma,\text{ans}}^{(k)}$ points to 'Local polynomial', $\mathcal{C}_{\gamma}^{(k)}$ points to 'Cut \rightarrow truth', and $\mathcal{R}_{\gamma,\text{MMC}}^{(k)}$ points to 'Rational'.

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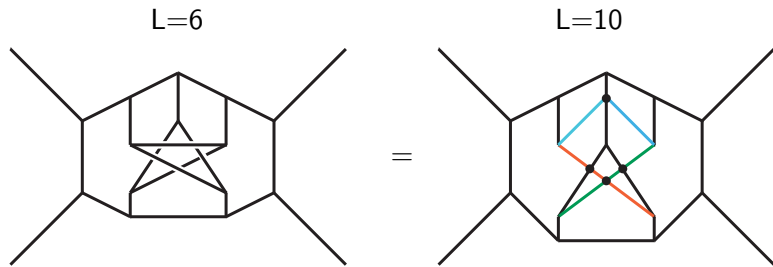
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$$\begin{array}{ccc}
 \text{Trying to determine} & \mathcal{P}_{\gamma, \text{ans}}^{(k)} = \mathcal{C}_{\gamma}^{(k)} - \mathcal{R}_{\gamma, \text{MMC}}^{(k)} & \text{Known higher cuts: } \sum \frac{n}{p^2} \\
 \downarrow & \downarrow & \downarrow \\
 & \text{Need to eval} &
 \end{array}$$

Applications of X ID in MMC

X Identity: Evaluate $\mathcal{C}^{(k)}$ *directly* from limit of **higher-loop** cut



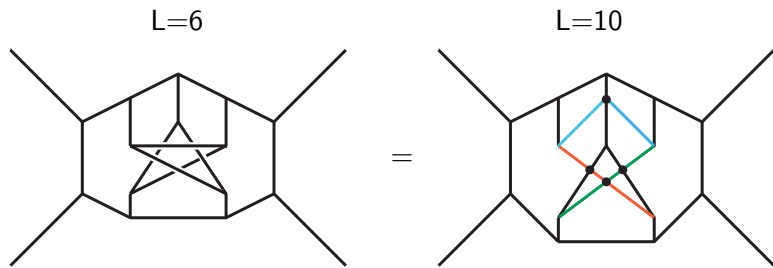
Challenges:

- ① Edge crossing is NP-Hard¹
- ② Quickly outpace known planar cuts (11+ loops)
- ③ Only works for color-ordered cuts

¹ex: There're better crossing schemes than in the diagram. Can you find one?

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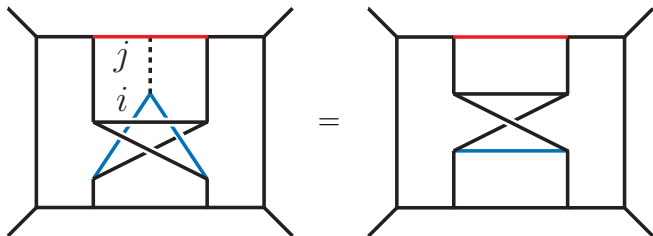
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$$\lim_{\ell_m \rightarrow 0} \mathcal{P}_{\gamma, \text{ans}}^{(k)} = - \left(\lim_{\ell_m \rightarrow 0} \mathcal{R}_{\gamma, \text{MMC}}^{(k)} \right) + (2p_i \cdot p_j) \mathcal{C}_{\gamma \setminus \ell_m}^{(k)}$$



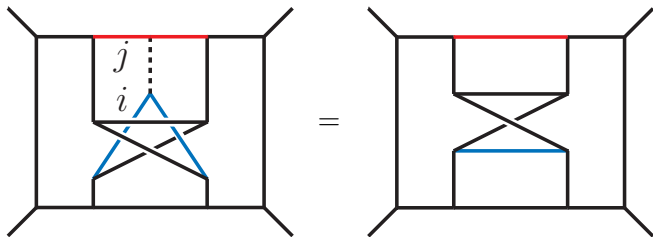
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- ① Need to merge conditions
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Challenges:

- 1 Need to merge conditions
- 2 Many evaluations of lower-loop cuts

Resolving the conditions

Fix a basis of momentum invariants

$\lim_{\ell_m \rightarrow 0}$ induces linear relations between invariants: π_{ℓ_m}

- Brute force: use an ansatz
 - \mathcal{P}_{ans} as literal polynomial ansatz in basis
 - H ID gives linear equations between ansatz parameters
- More clever: intersection of polynomial ideals

$$\pi_{\ell_m} \mathcal{P}_{\gamma, ans}^{(k)} = h_{\gamma, \ell_m} \mathcal{C}_{\gamma \setminus \ell_m}^{(k)} - \pi_{\ell_m} \mathcal{R}_{\gamma, MMC}^{(k)}$$

$$\Rightarrow \mathcal{P}_{\gamma, ans}^{(k)} = \pi_{\ell_m} \mathcal{P}_{\gamma, ans}^{(k)} + \ker \pi_{\ell_m} \subset \langle \pi_{\ell_m} \mathcal{P}_{\gamma, ans}^{(k)}, \pi_{\ell_m} \rangle$$

$$\mathcal{P}_{\gamma, ans}^{(k)} \sim \bigcap_{\ell_m \in \gamma} \langle \pi_{\ell_m} \mathcal{P}_{\gamma, ans}^{(k)}, \pi_{\ell_m} \rangle$$

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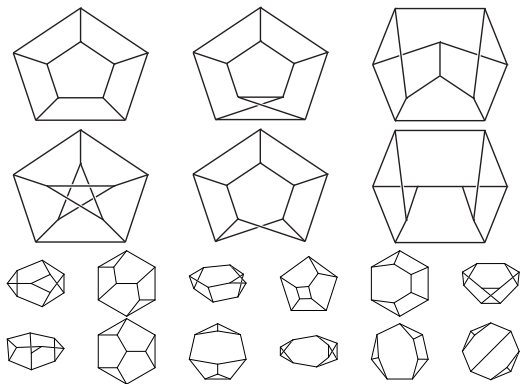
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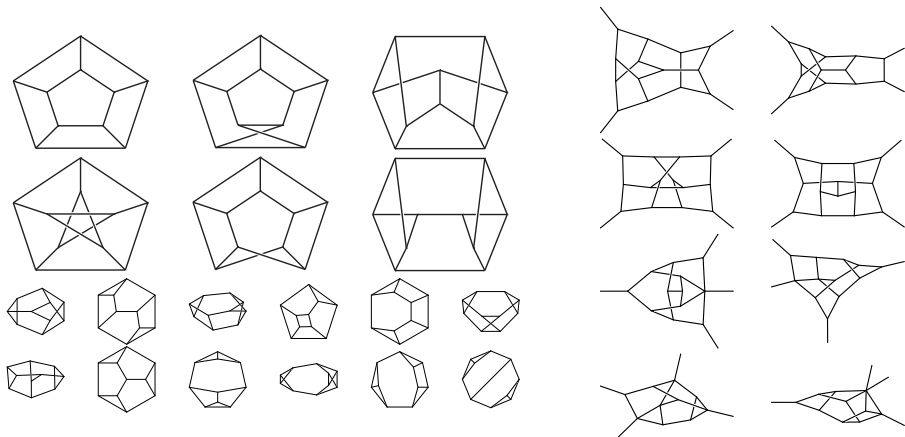
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Integrand Search Space



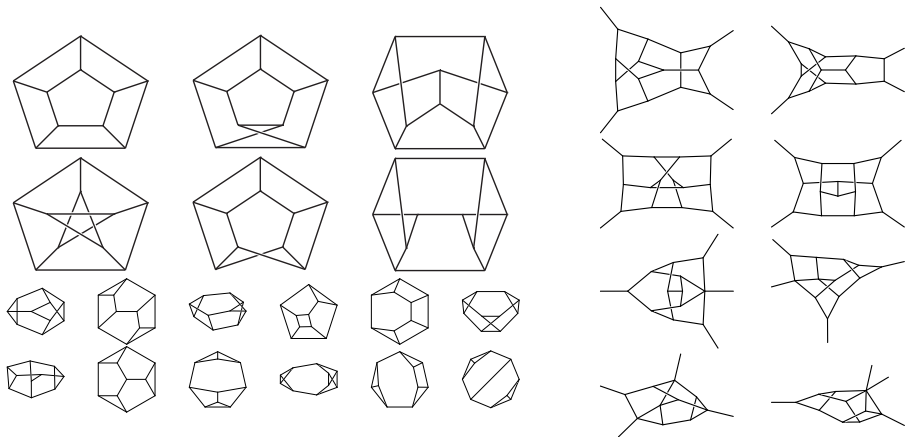
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# cuts	5548	41649	156853	363963	576582	706,281

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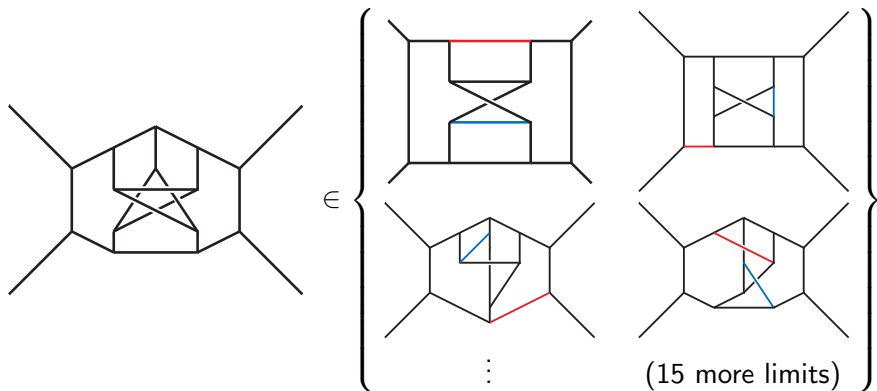
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6L Numerator Construction



The 6L Integrand

$N^k M$	0	1	2	3	4	Σ
cuts	5548	41649	156853	363963	576582	1.14×10^6
non-zero contacts	4420	16776	38000	54000	33000	1.46×10^5

Fits on a CD!

Contact terms carry half-ladder color factors

- Longest numerator: 62,511 terms
- Shortest numerator: 1 term (ladder diagrams)
- Average numerator: 90 terms

Looking Forward

- UV Integration
- Prepping tools for SUGRA: KLT, IBPs
- Improve efficiency for 7 loops
 - Intersection of ideals is sensitive to many superficial choices
 - Minimize number of limits to evaluate
 - Ansatz requires efficient inversion/row reduction
- Application to other theories: QCD, open string eff.
- Cubic representation: generalized double-copy, color-kinematics duality?

Thanks!

Questions?



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